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AUTHOR(S):

Murota, Kazuo

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# A Proof of the M-Convex Intersection Theorem

東京大学大学院 数理情報学専攻 室田 一雄 (Kazuo Murota)  
Graduate School of Information Science and Technology, University of Tokyo

## Abstract

This short note gives an alternative proof of the M-convex intersection theorem, which is one of the central results in discrete convex analysis. This note is intended to provide a direct simpler proof accessible to nonexperts.

## 1 M-Convex Intersection Theorem

The M-convex intersection theorem [3, Theorem 8.17] reads as follows, where  $V$  is a nonempty finite set, and  $\mathbf{Z}$  and  $\mathbf{R}$  are the sets of integers and reals, respectively; see §3 for the definitions of  $M^h$ -convex functions and notation  $\arg \min$ . This theorem is equivalent to the M-separation theorem, to the Fenchel-type min-max duality theorem, and to an optimality criterion of the M-convex submodular flow problem.

**Theorem 1 (M-convex intersection theorem).** *For  $M^h$ -convex functions  $f_1, f_2$  and a point  $x^* \in \text{dom} f_1 \cap \text{dom} f_2$  we have*

$$f_1(x^*) + f_2(x^*) \leq f_1(x) + f_2(x) \quad (\forall x \in \mathbf{Z}^V) \quad (1)$$

*if and only if there exists  $p^* \in \mathbf{R}^V$  such that<sup>1</sup>*

$$f_1[-p^*](x^*) \leq f_1[-p^*](x) \quad (\forall x \in \mathbf{Z}^V), \quad (2)$$

$$f_2[+p^*](x^*) \leq f_2[+p^*](x) \quad (\forall x \in \mathbf{Z}^V). \quad (3)$$

*For such  $p^*$  we have*

$$\arg \min(f_1 + f_2) = \arg \min f_1[-p^*] \cap \arg \min f_2[+p^*]. \quad (4)$$

*Moreover, if  $f_1$  and  $f_2$  are integer-valued, we can choose integer-valued  $p^* \in \mathbf{Z}^V$ .*

We shall give a constructive proof of Theorem 1 based on the successive shortest path algorithm. Different proofs available in [3] are:

1. original proof based on negative-cycle cancelling for the M-convex submodular flow problem (§9.5 and Note 9.21 of [3]), and
2. polyhedral proof for the discrete separation theorem based on the separation in convex analysis (Proof of Theorem 8.15 of [3]).

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<sup>1</sup>Notation:  $f_1[-p^*](x) = f_1(x) - \sum_{v \in V} p^*(v)x(v)$ ,  $f_2[+p^*](x) = f_2(x) + \sum_{v \in V} p^*(v)x(v)$ .

## 2 Essence of Theorem 1

The essence of Theorem 1 consists of two assertions:

1. optimality of  $x^* \Rightarrow$  existence of  $p^*$ ,
2. integrality of  $f_1, f_2 \Rightarrow$  integrality of  $p^*$ .

To see this we make easier observations in this section.

**Observation 1:** Existence of  $p^*$  with (2) and (3)  $\Rightarrow$  optimality (1) of  $x^*$ .

(Proof)

$$\begin{aligned} f_1(x^*) + f_2(x^*) &= f_1[-p^*](x^*) + f_2[+p^*](x^*) \\ &\leq f_1[-p^*](x) + f_2[+p^*](x) = f_1(x) + f_2(x). \end{aligned}$$

**Observation 2:** For any  $p^* \in \mathbf{R}^V$  we have

$$\arg \min(f_1 + f_2) \supseteq \arg \min f_1[-p^*] \cap \arg \min f_2[+p^*]. \quad (5)$$

(Proof) This follows from the inequality shown in the proof of Observation 1.

**Observation 3:** If

$$f_1[-p^*](x^\circ) \leq f_1[-p^*](x) \quad (\forall x \in \mathbf{Z}^V), \quad (6)$$

$$f_2[+p^*](x^\circ) \leq f_2[+p^*](x) \quad (\forall x \in \mathbf{Z}^V) \quad (7)$$

for some  $x^\circ$  and  $p^*$ , then

$$f_1[-p^*](x^*) \leq f_1[-p^*](x) \quad (\forall x \in \mathbf{Z}^V), \quad (8)$$

$$f_2[+p^*](x^*) \leq f_2[+p^*](x) \quad (\forall x \in \mathbf{Z}^V) \quad (9)$$

for every  $x^* \in \arg \min(f_1 + f_2)$ . Hence,

$$\arg \min(f_1 + f_2) \subseteq \arg \min f_1[-p^*] \cap \arg \min f_2[+p^*]. \quad (10)$$

(Proof) Put  $x = x^*$  in (6) and (7) to obtain

$$f_1[-p^*](x^\circ) \leq f_1[-p^*](x^*), \quad (11)$$

$$f_2[+p^*](x^\circ) \leq f_2[+p^*](x^*). \quad (12)$$

Adding these yields

$$\begin{aligned} f_1(x^\circ) + f_2(x^\circ) &= f_1[-p^*](x^\circ) + f_2[+p^*](x^\circ) \\ &\leq f_1[-p^*](x^*) + f_2[+p^*](x^*) = f_1(x^*) + f_2(x^*), \end{aligned}$$

whereas  $x^* \in \arg \min(f_1 + f_2)$ . Hence we have equalities in (11) and (12).

**Observation 4:** It suffices to consider M-convex functions rather than  $M^{\mathbf{I}}$ -convex functions.

(Proof) This follows from the equivalence between  $M^{\mathbf{I}}$ -convexity and M-convexity; see [3, §6.1].

Thus the proof of Theorem 1 is reduced to showing the following.

**Proposition 2.** For  $M$ -convex functions  $f_1, f_2$  with  $\arg \min(f_1 + f_2) \neq \emptyset$ , there exist  $x^\circ \in \arg \min(f_1 + f_2)$  and  $p^* \in \mathbf{R}^V$  such that

$$f_1[-p^*](x^\circ) \leq f_1[-p^*](x) \quad (\forall x \in \mathbf{Z}^V), \quad (13)$$

$$f_2[+p^*](x^\circ) \leq f_2[+p^*](x) \quad (\forall x \in \mathbf{Z}^V). \quad (14)$$

If  $f_1$  and  $f_2$  are integer-valued, we can choose integer-valued  $p^* \in \mathbf{Z}^V$ .

### 3 Notation and Basic Facts

We denote by  $\mathbf{Z}^V$  the set of integral vectors indexed by  $V$ , and by  $\mathbf{R}^V$  the set of real vectors indexed by  $V$ . For a vector  $x = (x(v) : v \in V) \in \mathbf{Z}^V$ , where  $x(v)$  is the  $v$ th component of  $x$ , we define the positive support  $\text{supp}^+(x)$  and the negative support  $\text{supp}^-(x)$  by

$$\text{supp}^+(x) = \{v \in V \mid x(v) > 0\}, \quad \text{supp}^-(x) = \{v \in V \mid x(v) < 0\}.$$

We use notation  $x(S) = \sum_{v \in S} x(v)$  for a subset  $S$  of  $V$ . For each  $S \subseteq V$ , we denote by  $\chi_S$  the characteristic vector of  $S$  defined by:  $\chi_S(v) = 1$  if  $v \in S$  and  $\chi_S(v) = 0$  otherwise, and write  $\chi_v$  for  $\chi_{\{v\}}$  for all  $v \in V$ . For a vector  $p = (p(v) : v \in V) \in \mathbf{R}^V$  and a function  $f : \mathbf{Z}^V \rightarrow \mathbf{R} \cup \{+\infty\}$ , we define functions  $\langle p, x \rangle$  and  $f[p](x)$  in  $x \in \mathbf{Z}^V$  by

$$\langle p, x \rangle = \sum_{v \in V} p(v)x(v), \quad f[p](x) = f(x) + \langle p, x \rangle.$$

We also denote the set of minimizers of  $f$  and the effective domain of  $f$  by

$$\arg \min f = \{x \in \mathbf{Z}^V \mid f(x) \leq f(y) \ (\forall y \in \mathbf{Z}^V)\},$$

$$\text{dom } f = \{x \in \mathbf{Z}^V \mid f(x) < +\infty\}.$$

A function  $f : \mathbf{Z}^V \rightarrow \mathbf{R} \cup \{+\infty\}$  with  $\text{dom } f \neq \emptyset$  is called  $M^h$ -convex if it satisfies

( $M^h$ -EXC) for all  $x, y \in \text{dom } f$  and all  $u \in \text{supp}^+(x - y)$ , there exists  $v \in \text{supp}^-(x - y) \cup \{0\}$  such that

$$f(x) + f(y) \geq f(x - \chi_u + \chi_v) + f(y + \chi_u - \chi_v),$$

where  $\chi_0$  is defined to be the zero vector in  $\mathbf{Z}^V$ .

A function  $f : \mathbf{Z}^V \rightarrow \mathbf{R} \cup \{+\infty\}$  with  $\text{dom } f \neq \emptyset$  is called  $M$ -convex if it satisfies

( $M$ -EXC) for all  $x, y \in \text{dom } f$  and all  $u \in \text{supp}^+(x - y)$ , there exists  $v \in \text{supp}^-(x - y)$  such that

$$f(x) + f(y) \geq f(x - \chi_u + \chi_v) + f(y + \chi_u - \chi_v).$$

A nonempty set  $B \subseteq \mathbf{Z}^V$  is called  $M$ -convex if it satisfies

(B-EXC) for all  $x, y \in B$  and all  $u \in \text{supp}^+(x - y)$ , there exists  $v \in \text{supp}^-(x - y)$  such that  $x - \chi_u + \chi_v, y + \chi_u - \chi_v \in B$ .

The minimizers of an M-convex function have a good characterization.

**Lemma 3** ([3, Theorem 6.26]). *For an M-convex function  $f$  and  $x \in \text{dom } f$ ,  $x \in \arg \min f$  if and only if  $f(x) \leq f(x - \chi_u + \chi_v)$  for all  $u, v \in V$ .*

**Lemma 4** ([3, Proposition 6.29]). *For an M-convex function  $f$ ,  $\arg \min f$  is an M-convex set if not empty.*

An M-convex set has the following property. (See [1, Lemma 4.5] and [2, Lemma 2.3.22, Remark 3.3.24]. This is a special case of [3, Proposition 9.23].)

**Lemma 5** (“no-short cut lemma”). *Let  $B$  be an M-convex set. For any  $x \in B$  and any distinct  $u_1, v_1, u_2, v_2, \dots, u_r, v_r \in V$ , if  $x - \chi_{u_i} + \chi_{v_i} \in B$  for all  $i = 1, \dots, r$  and  $x - \chi_{u_i} + \chi_{v_j} \notin B$  for all  $i, j$  with  $i < j$ , then  $y = x - \sum_{i=1}^r (\chi_{u_i} - \chi_{v_i}) \in B$ .*

## 4 Proof of Proposition 2 by SSP

We give a proof of Proposition 2 on the basis of the successive shortest path algorithm (SSP) [3, §10.3.4] as adapted to finding a minimizer of  $f_1 + f_2$ . We may assume that the effective domains of  $f_1$  and  $f_2$  are bounded.

Let  $x_1$  and  $x_2$  be arbitrary minimizers of  $f_1$  and  $f_2$ , respectively. We construct a directed graph  $G(f_1, f_2, x_1, x_2) = (V, A)$  and an arc length  $\ell \in \mathbf{R}^A$  as follows. Arc set  $A$  is the union of two disjoint parts:

$$\begin{aligned} A_1 &= \{(u, v) \mid u, v \in V, u \neq v, x_1 - \chi_u + \chi_v \in \text{dom } f_1\}, \\ A_2 &= \{(v, u) \mid u, v \in V, u \neq v, x_2 - \chi_u + \chi_v \in \text{dom } f_2\}, \end{aligned} \quad (15)$$

and  $\ell \in \mathbf{R}^A$  is defined by

$$\ell(a) = \begin{cases} f_1(x_1 - \chi_u + \chi_v) - f_1(x_1) & \text{if } a = (u, v) \in A_1, \\ f_2(x_2 - \chi_u + \chi_v) - f_2(x_2) & \text{if } a = (v, u) \in A_2. \end{cases} \quad (16)$$

The length function  $\ell$  is nonnegative due to Lemma 3.

Put  $S = \text{supp}^+(x_1 - x_2)$  and  $T = \text{supp}^-(x_1 - x_2)$ . A path exists from  $S$  to  $T$  by Lemma 6 below. Let  $P$  be a shortest path from  $S$  to  $T$  in  $G$  with a minimum number of arcs, and let  $t \in T$  be the terminal vertex of  $P$ .

Let  $d : V \rightarrow \mathbf{R} \cup \{+\infty\}$  denote the shortest distance from  $S$  to all vertices with respect to  $\ell$ . Then we have

$$\ell(a) + d(u) - d(v) \geq 0$$

for all arcs  $a = (u, v) \in A$ . Define  $p \in \mathbf{R}^V$  by  $p(v) = \min\{d(v), d(t)\}$  for all  $v \in V$ . It follows from the nonnegativity of  $\ell$  that

$$\ell(a) + p(u) - p(v) \geq 0$$

for all arcs  $a = (u, v) \in A$ . The above system of inequalities is equivalent to

$$\begin{aligned} f_1(x_1 - \chi_u + \chi_v) - f_1(x_1) + p(u) - p(v) &\geq 0, \\ f_2(x_2 - \chi_u + \chi_v) - f_2(x_2) - p(u) + p(v) &\geq 0 \end{aligned}$$

for all  $u, v \in V$ , which is further equivalent to

$$x_1 \in \arg \min f_1[-p], \quad x_2 \in \arg \min f_2[+p],$$

by Lemma 3. Note that for all arcs  $a = (u, v) \in A$ ,

$$\ell_p(a) = \ell(a) + p(u) - p(v)$$

are the lengths of  $a$  in the graph  $G(f_1[-p], f_2[+p], x_1, x_2)$  associated with  $f_1[-p]$ ,  $f_2[+p]$ ,  $x_1$ , and  $x_2$ .

Since  $\ell_p(a) = 0$  for all  $a \in P$ , we have

$$\begin{aligned} x_1 - \chi_u + \chi_v &\in \arg \min f_1[-p] && \text{for all } (u, v) \in P \cap A_1, \\ x_2 - \chi_u + \chi_v &\in \arg \min f_2[+p] && \text{for all } (v, u) \in P \cap A_2. \end{aligned} \quad (17)$$

Since  $P$  has a minimum number of arcs, we also have

$$x_1 - \chi_u + \chi_w \notin \arg \min f_1[-p], \quad x_2 - \chi_w + \chi_u \notin \arg \min f_2[+p] \quad (18)$$

for all vertices  $u$  and  $w$  of  $P$  such that  $(u, w) \notin P$  and  $u$  appears earlier than  $w$  in  $P$ .

Furthermore, arcs of  $A_1$  and  $A_2$  appear alternately in  $P$ . This can be proved as follows. Suppose that consecutive two arcs  $(u, v), (v, w) \in P$  belong to, say,  $A_1$ . Then, by (M-EXC),

$$f_1(x_1 + \chi_u - \chi_v) + f_1(x_1 + \chi_v - \chi_w) \geq f_1(x_1) + f_1(x_1 + \chi_u - \chi_w),$$

which yields

$$\ell(u, v) + \ell(v, w) \geq \ell(u, w),$$

a contradiction to the minimality (with respect to the number of arcs) of  $P$ . Consequently, we have

$$\begin{aligned} a_1 = (u_1, v_1), a_2 = (u_2, v_2) \in P \cap A_1, a_1 \neq a_2 &\implies \{u_1, v_1\} \cap \{u_2, v_2\} = \emptyset, \\ a_1 = (u_1, v_1), a_2 = (u_2, v_2) \in P \cap A_2, a_1 \neq a_2 &\implies \{u_1, v_1\} \cap \{u_2, v_2\} = \emptyset. \end{aligned} \quad (19)$$

From Lemmas 4 and 5 together with (17), (18), and (19), we have

$$x'_1 \equiv x_1 - \sum_{(u,v) \in P \cap A_1} (\chi_u - \chi_v) \in \arg \min f_1[-p], \quad (20)$$

$$x'_2 \equiv x_2 - \sum_{(v,u) \in P \cap A_2} (\chi_u - \chi_v) \in \arg \min f_2[+p]. \quad (21)$$

Thus the modification of  $(f_1, f_2, x_1, x_2)$  to  $(f'_1, f'_2, x'_1, x'_2)$ , where  $f'_1 = f_1[-p]$  and  $f'_2 = f_2[+p]$ , keeps the conditions

$$x'_1 \in \arg \min f'_1, \quad x'_2 \in \arg \min f'_2.$$

We have

$$x'_1 - x'_2 = (x_1 - x_2) - (\chi_s - \chi_t)$$

with  $s \in \text{supp}^+(x_1 - x_2)$  and  $t \in \text{supp}^-(x_1 - x_2)$ , since  $P$  is a path from  $\text{supp}^+(x_1 - x_2)$  to  $\text{supp}^-(x_1 - x_2)$  and arcs of  $A_1$  and  $A_2$  appear alternately in  $P$ . This implies that  $\sum_{v \in V} |x_1(v) - x_2(v)|$  is decreased by two. Repeating the modification above we eventually arrive at  $x_1 = x_2$ , when we have

$$x_1 \in \arg \min f_1[-p] \cap \arg \min f_2[+p].$$

Finally note that, if the functions  $f_1$  and  $f_2$  are integer-valued, the length function  $\ell$  is integer-valued, and hence  $p$  is also integer-valued.

The SSP algorithm is summarized below.

**Algorithm SSP** ( $f_1, f_2$ : M-convex)

**Step 0.** Find  $x_1 \in \arg \min f_1$  and  $x_2 \in \arg \min f_2$ . Set  $p := 0$ .

**Step 1.** If  $x_1 = x_2$  then stop.

**Step 2.** Construct  $G$  and compute  $\ell$  for  $f_1[-p]$ ,  $f_2[+p]$ ,  $x_1$  and  $x_2$  by (15) and (16).  
Set  $S := \text{supp}^+(x_1 - x_2)$ ,  $T := \text{supp}^-(x_1 - x_2)$ .

**Step 3.** Compute the shortest distances  $d(v)$  from  $S$  to all  $v \in V$  in  $G$  with respect to  $\ell$ .  
Find a shortest path  $P$  from  $S$  to  $T$  with a minimum number of arcs, and let  $t$  be the terminal vertex of  $P$ .

**Step 4.** For all  $v \in V$ , set  $p(v) := p(v) + \min\{d(v), d(t)\}$ .  
Update  $x_1$  and  $x_2$  by (20) and (21).  
Go to Step 1.

**Lemma 6.** If  $\text{dom } f_1 \cap \text{dom } f_2 \neq \emptyset$  and  $x_1 \neq x_2$ , then there exists a path from  $S = \text{supp}^+(x_1 - x_2)$  to  $T = \text{supp}^-(x_1 - x_2)$ .

*Proof:* To prove by contradiction, suppose that there exists no path from  $S$  to  $T$  and let  $W$  be the set of the vertices reachable from  $S$ . Then  $W \supseteq S$  and  $W \cap T = \emptyset$ .

Define set functions  $\rho_i : 2^V \rightarrow \mathbf{Z} \cup \{+\infty\}$  as

$$\rho_i(X) = \sup\{z(X) \mid z \in \text{dom } f_i\}$$

for  $i = 1, 2$ . For  $z \in \text{dom } f_i$  we obviously have<sup>2</sup>

$$z(X) \leq \rho_i(X) \quad (\forall X \subseteq V).$$

<sup>2</sup>As is well known (see [3, §4.4]), the M-convexity of  $\text{dom } f_i$  implies that  $\rho_i$  is submodular and

$$\text{dom } f_i = \{z \in \mathbf{Z}^V \mid z(X) \leq \rho_i(X) \ (\forall X \subset V), z(V) = \rho_i(V)\}.$$

However, we do not need this fact for the proof of Lemma 6.

We also have  $z(V) = \rho_i(V)$  since  $y(V)$  is constant for all  $y \in \text{dom } f_i$ . Hence, for all  $z \in \text{dom } f_1 \cap \text{dom } f_2$  we have

$$\rho_1(V) = z(V) = z(V \setminus X) + z(X) \leq \rho_1(V \setminus X) + \rho_2(X) \quad (\forall X \subseteq V). \quad (22)$$

Since  $x_1 \in \text{dom } f_1$  and there exists no arc of  $A_1$  from  $W$  to  $V \setminus W$ , we have

$$x_1(V \setminus W) = \rho_1(V \setminus W)$$

by Lemma 3 applied to an M-convex function

$$f(z) = \begin{cases} -z(V \setminus W) & \text{if } z \in \text{dom } f_1, \\ +\infty & \text{otherwise.} \end{cases}$$

Symmetrically, since  $x_2 \in \text{dom } f_2$  and there exists no arc of  $A_2$  from  $W$  to  $V \setminus W$ , we have

$$x_2(W) = \rho_2(W).$$

Adding these yields

$$x_1(V) - [x_1(W) - x_2(W)] = \rho_1(V \setminus W) + \rho_2(W).$$

This contradicts (22), since  $x_1(V) = \rho_1(V)$  and  $[x_1(W) - x_2(W)] > 0$  by  $W \supseteq S$  and  $W \cap T = \emptyset$ .

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